Reminders:
1. Next exam is a take-home; handed out 3/17/00, due 3/20/00 (5pm).
2. Exam will cover through Lecture 18 (on the web site). You are responsible for the text, through Intermezzo 3 (on local) and for the contents of the lab lectures up to lab 6 (on the web site).
3. Homework available this afternoon; it will be 1/2 normal size.
4. Go to the lab lecture this week. Many functions that will be needed in future labs and lectures will only be covered in the lab lecture.
5. Should we do a review session? (At most one session)

Review
1. We started to talk about functional abstraction. Built a series of repetitive examples, intended to show you that they were similar. Showed how to abstract them out. Ended with an example, keep-rel: (num num -> num) num alon -> alon.

```
;; keep-rel (num num -> num) num list-of-ums -> list-of-ums
;; Purpose: keep all the numbers in the input list that have the relation given
;; by the function argument to the number argument (whew!)
(define (keep-rel relation num alon)
  (local [(define filter-rel alon) ;; treat relation & num as invariant
           (cond
             [(empty? alon) empty]
             [(cons? alon)
              (cond
                [(relation (first alon) num)
                 (cons (first alon) (filter-rel (rest alon)))]
                [else (filter-rel (rest alon))] ) ]))
        (filter-rel alon))

(define (keep-gt-9 alon)
  (keep-rel > 9 alon))
```

This shocked several of you. We were able to pass the relational operator as an argument to a function and it works.

Programs as Values
How can we pass > or < or = as an argument (parameter) to a function such as keep-rel? A program is just another value. Recall that we created test values using the define operation. We used it to create numbers and lists. These became named values in the world of Scheme values. We use the same syntax and operation to create programs, don’t we? The define operator takes its two parts – the function name and argument list and its body. The way that DrScheme evaluates the function is a little more complex than evaluating (define Seventeen 17) because it involves rewriting the parameters with their
values. Of course, Seventeen has no parameters, so the mechanism for handling functions will work on a defined numerical value as a trivial case.

In general, functions (programs) are values. We can use them anywhere that we use values. For example, we can build lists of functions, just as we build lists of symbols or numbers, or objects created with define-struct.

What about > ? Isn't it special? Isn't it a built-in operator?

Yes. It is a built-in operator, but that grants it no special exemptions from the rules that govern the execution of Scheme programs. The built-in operators are just functions themselves. Thus, +, *, /, -, <, =, > are all programs that you could write. So are add1, sub1. BUT, define and define-struct are not implemented that way; they are only valid in the definitions window—not in the interactions window. That gives you a hint that they are not Scheme operators.

Back to Work
So far, all of these examples have looked at an open-ended interval (> 5, < 9). [Of course, we could use filter-rel to find all of the numbers equal to x, but that's only interesting if we wanted to count them. (length (keep-rel = 5 somelist)).] What if we wanted to pull out the numbers between five and nine?

;; keep-bet-5-9: list-of-numbers -> list-of-numbers
;; Purpose: returns a list containing those numbers in the input list
;; whose value is between 5 and 9, inclusive
(define (keep-bet-5-9 alon)
  (cond
   [(empty? alon) empty]
   [(cons? alon)
    (cond
     [(and (>= (first alon) 5) (<= (first alon) 9))
      (cons (first alon) (keep-bet-5-9 (rest alon)))]
     [else (keep-bet-5-9 (rest alon))]
    ) ] ) )

We should really write a helper function to replace the complex test.

;; bet-5-9?: number -> boolean
;; Purpose: test if the argument is between five and nine, inclusive
(define (bet-5-9? anum)
  (and (>= num 5) (<= num 9)))
(define (keep-bet-5-9 alon)
  (cond
   [(empty? alon)   empty]
   [(cons?    alon)
     (cond
      [(bet-5-9? (first alon))
       (cons (first alon) (keep-bet-5-9 (rest alon)))]
      [else  (keep-bet-5-9 (rest alon))
       ) ] ) )
)

You know, by now, where we are going. What if we want to change the range of numbers? We can change the helper function, but we might be using it somewhere else. We can write a version that takes an arbitrary range of numbers…

;; bet? :  num num num -> boolean
;; Purpose: determines if the third argument lies numerically between the
;;          first and second arguments
(define (bet?  lower upper  anum)
  (and  (>= num lower) (<= num upper)))

;; keep-bet : num num list-of-numbers -> list-of-numbers
;; Purpose: keeps all the numbers lying between first and second arguments
(define (keep-bet  lower  upper  alon)
  (local
   [(define (filter-bet alon)
     (cond
      [(empty? alon)  empty]
      [(cons?    alon)
        (cond
         [(bet? lower upper (first alon))
          (cons (first alon) (filter-bet (rest alon)))]
         [else  (filter-bet  (rest alon))
          ) ] ) ] )

  (filter-bet alon) ))

(define (keep-bet-5-9 alon)
  (keep-bet 5 9 alon))

Notice that we used a local to avoid passing around lower and upper at the recursive calls inside filter-bet and keep-bet.

**Abstracting from keep-rel & keep-bet**

Look at the definitions of keep-rel and keep-bet. They are reasonably similar. They differ in their parameters and the first case in the innermost cond—where they decide
whether or not to keep the first element of the input list. The parameter differences boil down to inputs to that test. The cond clause differs in the implementation of that test.

Can we write one program to capture all of that common code? We start by copying down all of the information that is common to both programs, leaving an ellipsis in places where they differ…

\[
\text{(define (keep keep-elt? alon)}
\text{(local}
\text{[(define (filter alon)}
\text{(cond}
\text{[(empty? alon) empty]}}
\text{[(cons? alon)}
\text{(cond}
\text{[(... (first alon))}
\text{(cons (first alon) (filter (rest alon))))}
\text{[else (filter (rest alon))]})])}
\text{(filter alon))])}
\]

Let's fill in the gaps. First, look at the gap inside \textit{filter}. We need a helper function that takes (first alon) and returns a boolean that tells \textit{filter} whether or not to keep the number. Let's make a number for that helper function and make it a parameter. That fills in the gap in the parameters, too.

\[
\text{(define (keep keep-elt? alon)}
\text{(local}
\text{[(define (filter alon)}
\text{(cond}
\text{[(empty? alon) empty]}}
\text{[(cons? alon)}
\text{(cond}
\text{[(keep-elt? (first alon))}
\text{(cons (first alon) (filter (rest alon))))}
\text{[else (filter (rest alon))]})])}
\text{(filter alon))})}
\]

To use this, we just need to write the appropriate helper functions. For example

\[
\text{(define (keep-lt-5 alon)}
\text{(local [(define (lt-5? num) (< num 5))]
\text{(keep lt-5? alon))})}
\]

\[
\text{(define (keep-bet-5-9 alon)}
\text{(local [(define (bet-5-9? num) (bet? 5 9 num))]
\text{(keep bet-5-9? alon))})}
\]
The function **keep** is so useful that Scheme provides a built-in version of it. We call the built-in version of it **filter**.

**The Professor is a Liar**

Well, I told you that everything you learned outside of the compilers class is a lie. In fact, Scheme does provide **filter**, but it is not as limited as the one we wrote.

What if you wanted to write a function that counted the number of times the symbol ‘fee appeared in a list? You'd like to write it as

```scheme
;; keep-fee : list-of-symbol -> list-of-symbol
;; Purpose: return the list containing every occurrence of 'fee
(define (keep-fee alos)
  (local [(define (is-fee? asym)(= 'fee asym))]
    (keep is-fee? alos)))
```

Except, this violates the contract for keep.

Look at the body of keep. Is there anything in the body of keep that depends on the fact that the argument list contains numbers?

The contract is overly restrictive. In fact, we'd like keep to work on a list of symbol, or a list of ftn, or a list of mechanic, or a list of plane, or a list of just-about-anything.

To do this, we could rewrite the contract by defining a list of several types of things. Imagine **list-of-just-about-everything** where it has a clause for each kind of object that we want keep to use. Does this solve our problem? NO. In particular, the list-of-just-about-everything allows one list to contain many different kinds of elements. That makes the keep-elt function much harder to write correctly. What we need is a notation that allows us to specify that keep takes a list of something—where every element of the list is a something—but that keep doesn't care what that something is. The precise requirement for keep is that keep-elt must have a contract of **something-> boolean**.

To write this down as a Scheme contract, we want to say

```scheme
;; a list-of-alpha is either
;;   – empty, or
;;   – (cons f r) where f is an alpha and r is a list of alpha.
```

Then, we can write the contract for keep as

```scheme
;; keep : (alpha->boolean) list-of-alpha -> list-of-alpha
```

This concisely expresses the requirement for keep. The function keep-elt must process the elements of the argument list and return a boolean. The elements of the input and output list both have that same type. This is the real contract for filter.