Exam

- Solutions will be posted today or tomorrow
- Look at the solutions

Homework 9

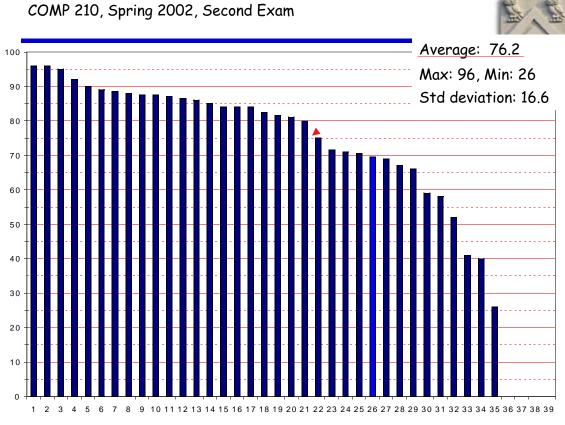
(*Ex. 32.2.1 – 32.2.8 in book*)

- Due Wednesday, April 10, 2002 in class
- Do one sub-problem each day and you will finish early
- Procrastinate and you will <u>not</u> finish

Labs this week as normal

Challenge lab?

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Finishing up accumulators

The example with reverse was tortured (my fault)

• Can we write another classic program with an accumulator?

 \rightarrow Let's try max, one of our favorite examples

;; a non-empty-list-of-number (nelon) is either

;; - (cons f r) where f is a number and r is empty, or

;; $-(\cos f r)$ where f is a number and r is a nelon

;; We will use Scheme's built-in list constructor to implement nelons

;; maxacc: nelon -> number ;; Purpose: returns the largest entry in a non-empty list of numbers (define (maxacc anelon) ...)

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Finishing up accumulators

Max, again

;; maxacc: nelon -> number

;; Purpose: returns the largest entry in a non-empty list of numbers (define (maxacc anelon) ...)

How do we proceed?

- With an accumulator, can pass along largest element so far
- What does helper do?

;; maxh: nelon number -> number ;; Purpose: returns the larger of acc and (max-of-list anelon) ;; acc holds the largest element seen so far (define (maxh anelon acc) ...)







Finishing up accumulators

Focusing on maxh

;; maxh: nelon number -> number ;; Purpose: returns the larger of acc and (max-of-list anelon) (define (maxh anelon acc) (cond [(empty? anelon) acc] [(> (first anelon) acc) [(else

```
))
```

(maxh (rest anelon) (first anelon))] (maxh (rest anelon) acc)]

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Focusing on maxh

;; maxh: nelon number -> number ;; Purpose: returns the larger of acc and (max-of-list anelon)

(define (maxacc anelon acc)

(cond

```
[(empty? anelon)
[(> (first anelon) acc)
[(else
```

accl (maxh (rest anelon) (first anelon))] (maxh (rest anelon) acc)]

But wait

))

- Maxh tests (empty? anelon)
- How can a nelon be empty?
- We subtly changed the problem & the contract •





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Maxh operates on a list

;; maxh: alon number -> number ;; Purpose: returns the larger of acc and (max-of-list alon) (define (maxh alon acc) (cond [(empty? alon) acc] [(cons? alon) (cond [(> (first alon) acc) (maxh (rest alon) (first alon))] [else (maxh (rest alon) acc)])]

- maxacc takes a nelon & uses (first anelon) as initial accum'r
- <u>maxh</u> takes a list & returns a number
 - \rightarrow Uses (empty? alon) test to return accumulator value

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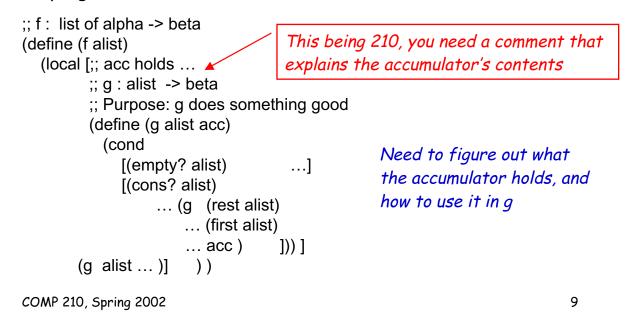
Finishing up accumulators

Putting it together

```
;; maxacc: nelon -> number
;; Purpose: returns the largest entry in a non-empty list of numbers
(define (maxacc anelon)
 (cond
    [(empty? (rest anelon)) (first anelon)]
    (cons? (rest anelon))
     (local
       ;; maxh: alon number -> number
         ;; Purpose: returns the larger of acc and (max-of-list anelon)
         (define (maxh alon acc)
           (cond
             [(empty? alon)
                                              accl
             [(> (first alon) acc)
                                              (maxh (rest alon) (first alon))]
             [(else
                                              (maxh (rest alon) acc)]))]
       (maxh (rest anelon) (first anelon)) )]
   ))
```

An aside

• We can think of this example as a template for accumulator programs over lists



Finishing up accumulators

What's the point?

- Old version of max worked
 - \rightarrow Used local to make it run in linear time (rather than 2^N)

```
;; maxclassic: nelon -> number
;; Purpose: rehash max, again
(define (maxclassic anelon)
  (cond
    [(empty? (rest anelon)) (first anelon)]
    [(cons? (rest anelon))
    (local [(define maxrest (maxclassic (rest anelon)))]
        (cond
        [(> (first anelon) maxrest) (first anelon)]
        [else maxrest]
        ))]
    ))
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```



What's the point?

- Old version of max worked
 - \rightarrow Used local to make it run in linear time (rather than 2^N)
- Does <u>maxacc</u> differ from <u>maxclassic</u> in any useful way
 - \rightarrow Consider their behavior on (list 1 2 3 4)

This is a point I tried to make with reverse last class Using the stepper made it particularly hard to see the point

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Finishing up accumulators

Consider the evaluation of each function (maxclassic (list 1 2 3 4))

⇒ defines maxrest0 as (maxclassic (list 2 3 4)
 ⇒ defines maxrest1 as (maxclassic (list 3 4))
 ⇒ defines maxrest2 as (maxclassic (list 4))
 ⇒ This returns 4
 ⇒ evaluates the cond and returns 4
 ⇒ evaluates the cond and returns 4
 ⇒ evaluates the cond and returns 4





Consider the evaluation of each function (maxacc (list 1 2 3 4)) \Rightarrow finds (rest anelon) is non-empty & enters local \Rightarrow evaluates (maxh (list 2 3 4) 1)

 \Rightarrow evaluates (maxh (list 3 4) 2) \Rightarrow evaluates (maxh (list 4) 3) \Rightarrow evaluates (maxh empty 4) & returns 4 \Rightarrow returns 4 \Rightarrow returns 4 \Rightarrow returns 4

What's the difference?

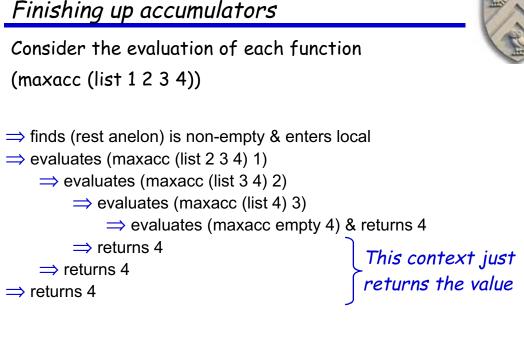
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```
Finishing up accumulators
Consider the evaluation of each function
(maxclassic (list 1 2 3 4))
\Rightarrow defines maxrest0 as (maxclassic (list 2 3 4)
     \Rightarrow defines maxrest1 as (maxclassic (list 3 4))
          \Rightarrow defines maxrest2 as (maxclassic (list 4))
               \RightarrowThis returns 4
          \Rightarrow evaluates the cond and returns 4
                                                       This context involves
further evaluation
     \Rightarrow evaluates the cond and returns 4
\Rightarrow evaluates the cond and returns 4
```

Scheme has lots of pending context after the recursive call







Scheme has (almost) no context after the call

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Does this matter?

- In large evaluations, that extra context adds up
- Takes space (in DrScheme) and time
- Can become a source of inefficiency

Tail recursion

- A tail-recursion returns the value of a self-recursive call
 - $\rightarrow~$ No further computation
- This is a particularly efficient form of recursion
 - \rightarrow Most translators (like DrScheme) optimize for this case





Another use for accumulators

- We can use an accumulator to transform a program into tailrecursive form
- This is an efficiency hack
 - \rightarrow But can be an important one

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