## Administrative Announcements

- Homework due today
- Next homework available today, due next Friday
- Challenge lab: tonight at 8:30 in Ryon
- Exam next Wednesday night
$\rightarrow$ Covers lecture through Friday, lab lectures
$\rightarrow 7$ to 9 pm
$\rightarrow$ Closed notes, closed book
$\rightarrow$ Location TBA
$\rightarrow$ Wednesday night lab folks should attend another lab


## Review

Last lecture:

- Did a whole series of examples
$\rightarrow$ keep-It-x, keep-gt-y, keep-bet-u-and-v
- Used parameterization to share code
- Used local to simplify the code
- Finally, abstracted out the conceptual heart of the code filter: (alpha->boolean) list-of-alpha -> list-of-alpha
$\rightarrow$ We call filter an abstract function
$\rightarrow$ We will encounter more abstract functions
$\rightarrow$ We will make heavy use of them


## Review

Develop keep-fee
;; keep-fee: list-of-symbol -> list-of-symbol
;; Purpose: returns a list containing every occurrence of 'fee' in the list
;; (define (keep-fee alos) ...)
(keep-fee (list 'fee 'fie 'foe 'fum 'fee)) -> (list 'fee 'fee)
(keep-fee empty) -> empty
(define (keep-fee alos)
(local [(define is-fee? asym) (symbol=? asym 'fee))]
(filter is-fee? alos)
))

COMP 210, Spring 2002

## Review

Critical points

- Pass a program as an argument
$\rightarrow$ Description is its contract in parentheses
$\rightarrow$ cons would be (alpha list-of-alpha $->$ list-of-alpha)
- Scheme functions are just programs*
$\rightarrow$ Can pass cons, $\langle\rangle,+,$, symbol $=$ ? as arguments
$\rightarrow$ Programs are data

This is not your basic
high-school AP
programming course

- Concept is called functional abstraction


## Helper functions

Abstract functions usually require helper functions

- Create many new names
$\rightarrow$ Cognitive overhead of inventing and tracking names
$\rightarrow$ Helper functions are used once, as was is-fee?
- Can hide them inside a local
$\rightarrow$ Works fine
$\rightarrow$ Well-understood rewriting rules
- But, ...
$\rightarrow$ A fairly heavy price to pay for creating and using a function
$\rightarrow$ Lots of typing, lots of steps in rewriting rules


## Local for helper functions

Using local for this purpose is hard to justify


- Our rules for local

Elliding invariants fits either one

1. Use local to avoid computing some complicated value more than once. This made a huge difference in the cost of max.
2. Use local to make complex expressions more readable by introducing helper functions that break it into tractable parts.

- This case doesn't really fit either criterion
$\rightarrow$ The expression is used once, not twice, or thrice, or ...
$\rightarrow$ The expression is not complicated.
$\rightarrow$ is-fee? is about as simple as Scheme gets ...
- We used a local just to create a function that we can pass to filter


## Helper functions

Need the ability to create anonymous functions

- Want a quick, easy, compact syntax
- Should create full-fledged functions

Enter $\lambda$, written lambda

- Lambda is a constructor for anonymous functions
(lambda $\left(\arg _{1} \arg _{2} \ldots \arg _{n}\right)$ expression)
$\rightarrow$ Creates an anonymous function of $n$ arguments
$\underset{\text { (symbol=? asym 'fee)) }}{\text { (define (is-fee? asym) }} \equiv \quad \equiv \quad \begin{gathered}\text { (lambda (asym) } \\ \text { (symbol=? asym 'fee)) }\end{gathered}$

COMP 210, Spring 2002

## Using lambda

We can use an anonymous function in keep-fee
;; keep-fee: list-of-symbol -> list-of-symbol
;; Purpose: return a list containing each occurrence of 'fee (define (keep-fee alos)
(filter (lambda (asym)(symbol=? asym 'fee)) alos))

This is equivalent to our earlier version of keep-fee
;; keep-fee: list-of-symbol -> list-of-symbol
;; Purpose: return a list containing each occurrence of 'fee (define (keep-fee alos)
(local [(define is-fee? asym) (symbol=? asym 'fee))]
(filter is-fee? alos)
))

## Using lambda

What does lambda do?
Dr. Scheme rewrites (lambda $\left(\arg _{1} \arg _{2} \ldots \arg _{n}\right)$ expression) as

```
(local [(define (a-unique-name arg}\mp@subsup{\mp@code{Arg}}{2}{}\ldots..\mp@subsup{\operatorname{arg}}{n}{}\mathrm{ )
        expression)
]
a-unique-name)
```

Subtle points

- The rewriting process has to concoct the name, not you
- This creates the function \& returns it


## Another example

Develop squares
;; squares: list-of-number -> list-of-number
;; Purpose: returns a list containing the squares of the input list


It would be cleaner to use a helper function, square

## Another example

## Develop squares

;; squares: list-of-number -> list-of-number
;; Purpose: returns a list containing the squares of the input list (define (squares alon)
(local [(define (square x)(* x x))]
(cond
[(empty? alon) empty]
[(cons? alon) (cons (square (first alon)) (squares (rest alon)) )] )))

We could develop cubes, \& quads, \& quints, \& ...

- These need helper functions cube, quad, quint, ...
- They fit a pattern: apply function to every element of a list


## Another abstract function

Scheme provides the abstract function map

- Takes function \& list
- Applies function to list, element-by-element
;; squares: list-of-number -> list-of-number
;; Purpose: returns a list containing the squares of the input list (define (squares alon)
(map (lambda (x)(* x x)) alon))

