

From Programs to Executions: An Odyssey in Language Translation

(with examples in Scheme)

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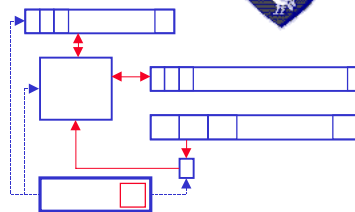


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What commands does the “computer” run?

Computer’s *instruction set*

- Low-level, imperative commands
 - > Arithmetic operations
 - > Memory operations
 - > Control operations
 - > Location-oriented programming
- We call these operations “assembly-language”



<u>Arithmetic Operations</u>
add r1, r2 => r3
sub r1, r2 => r3
mult r1, r2 => r3
div r1, r2 => r3

<u>Memory Operations</u>
load r1 => r2
store r1 => r2
loadi c1 => r2
copy r1 => r2

<u>Control Operations</u>
branch r1 -> r2
branchi r1 -> L2
call -> L1
return

Programming with Machine Operations



```
(define (summation n)
  (cond [(= n 0) 0]
        [else (+ n
                  (summation (sub1 n)))]))
```

This might become, after storage
assignment & translation

```
100: n
101: result
```

Assembly programming &
assemblers in COMP/ELEC 320

```
0: loadi 1    => r0
1: loadi 100 => r1
2: eq?  r1 r1 => r11
3: load  r1    => r2
4: copy  r0    => r3
5: add   r3 r2 => r3
6: sub   r2 r0 => r2
7: eq?   r2 r0 => r10
8: branch r10  -> 10
9: branch r11  -> 5
10: loadi 101 => r11
11: store r2   => r11
12: stop
```

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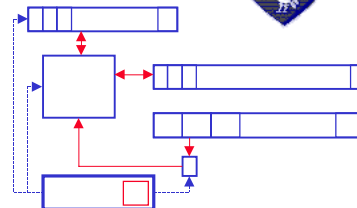
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Understanding How the Computer Works



One valuable tool is simulation

- Write a program that has the same behavior
- Models behavior of system
- Gives insight into its workings



Simulation is used in many ways

- Design of new systems
- Conduct experiments that are expensive or dangerous
- Train pilots in cases where loss is expensive

Simulating a computer shows us a lot about how it works

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How Does the Computer Work?



```
;; an op is a structure
;; (make-op fn arg1 arg2 arg3)
(define-struct op (fn arg1 arg2 arg3))

;; program memory holds 10k operations
(define Prog (build-vector 10240
                          (lambda (x) (make-op 'STOP 0 0 0))))

;; data memory holds 10k "word" of data
(define Data (build-vector 10240
                          (lambda (x) 0)))

;; the program counter holds the address of the current op
(define PC 0)
```

How Does the Control Unit Work?



```
;; control-unit : void -> boolean
;; Purpose: execute the operations in program memory, in
;; sequential order
(define (ControlUnit)
  (local [(define ThisOp (vector-ref Prog PC))]
    (begin
      (set! PC (add1 PC)) ;; update next operation's position
      (cond [(symbol=? (op-fn ThisOp) 'STOP) true]
            [else
             (begin
                (FunctionalUnit ThisOp)
                (ControlUnit)) ]])
      )))
```

How Does the Functional Unit Work?



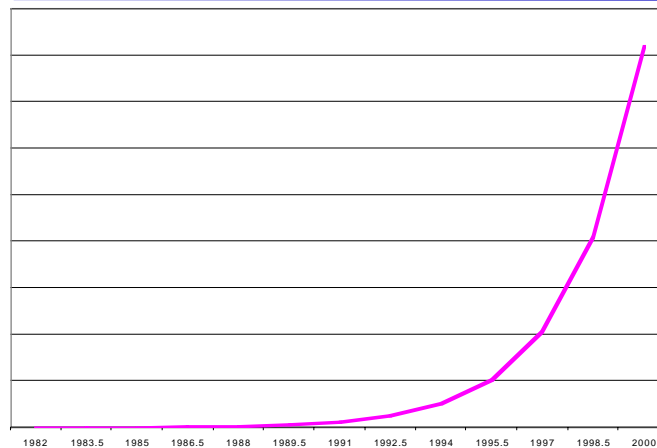
```
;; FunctionalUnit : op -> void
(define (FunctionalUnit AnOp)
  (local [ (define f (op-fn AnOp)) (define r1 (op-arg1 AnOp))
          (define r2 (op-arg2 AnOp)) (define r3 (op-arg3 AnOp)) ]
    (cond
      [(symbol=? f 'ADD)
       (vector-set! Regs r3 (+ (vector-ref Regs r1) (vector-ref Regs r2)))]
      [(symbol=? f 'LOAD)
       (vector-set! Regs r2 (vector-ref Data (vector-ref Regs r1)))]
      ...
      [(symbol=? f 'EQ?)
       (vector-set! Regs r3 (= (vector-ref Regs r1)(vector-ref Regs r2)))]
      [(symbol=? f 'BRANCH)
       (cond
         [(vector-ref Regs r1) (set! PC r2)] ;; not Regs[r2], but r2 itself
         [else void])]
      )))
```

[Complete code on web site](#)

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Computers Keep Getting Faster



Plot of
Moore's Law
1980-2000

How did this
happen?

Processor power versus time, 1980 to 2000

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They Are Running Faster



The clock frequency of processors has risen

- 1983: 10 MHz 68020 provided about 1 MIP
- 2000: 500 MHz PowerPC provides 500 MIPS + 1 GFLOP
- 2000: Pentiums in 700 to 800 MHz range

⇒ 10 cycles/operation

High-end chips are breaking the 1 GHz barrier

⇒ 1 cycle/operation

All this power has a downside, however

- Power consumption \propto frequency²
- Heat \propto power
- Computation needs operands, needs memory

They Are Also Running More Operations



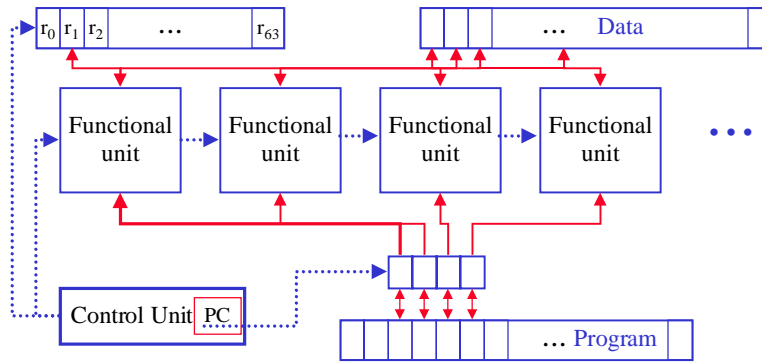
Programs contain parallelism

- Operations that can execute at the same time
- Can occur at the fine-grained level or on a larger scale

Computers can exploit parallelism

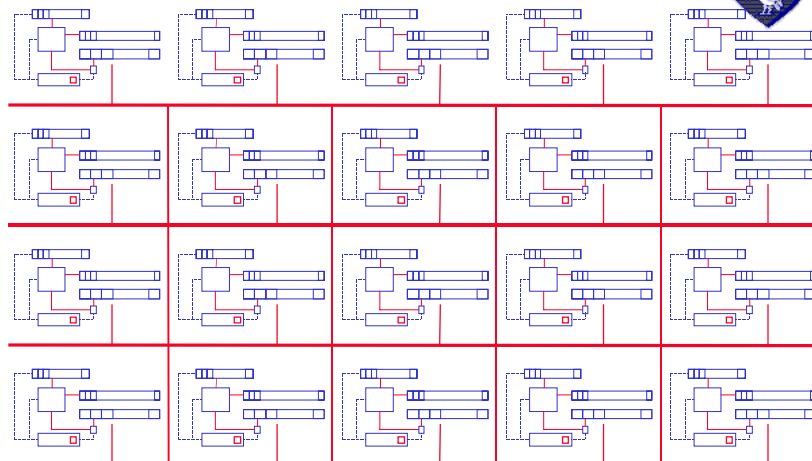
- Increase operations per cycle
- Use more hardware rather than faster hardware
- Many options
 - > Single chip processor with many functional units
 - > Custom-built machine with many individual nodes (computers)
 - > Networks of workstations and/or PCs

Single Chip, Many Functional Units



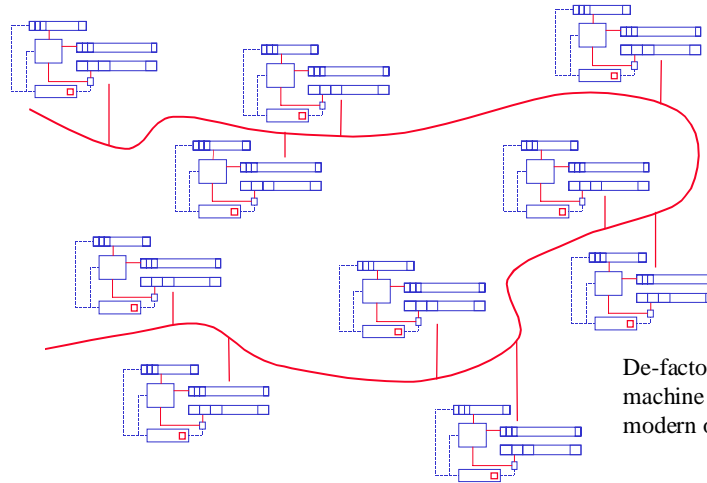
This takes advantage of instruction-level parallelism

Many Processors, Dedicated Interconnect



Multiprocessor Computer

Network of Workstations



De-facto parallel machine exists in most modern offices !

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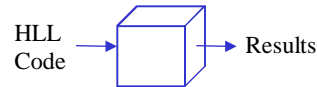
Going From HLL to ASM

Assembly programming & assemblers in COMP/ELEC 320

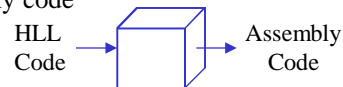


How do we get from a high-level language to assembly code for some target machine?

- Write an interpreter in a working HLL
 - > The base of DrScheme is written in C & C++
 - > The rest is written in Scheme
 - > Subject of COMP 311



- Write a compiler for the HLL
 - > Translate HLL directly to assembly code
 - > Must run the assembly code
 - > Subject of COMP 412



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Inside a Compiler

Chooses names, operations, maps
n-ary operators into binary ops.



- Represent the program in some internal form
(+ a b c) ⇒ (list + a b c)
- Traverse that data structure and generate code

(list + a b c) ⇒

```
add a,b ⇒ t1
add t1,c ⇒ t2
print t 1
```

Along the way

- Do most of the things that the interpreter does!
- Plan, when compiler runs, how events will occur when the program actually runs
- Plan once, execute many times

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Compiler versus Interpreter



- The interpreter manages resources as demands arise
- The compiler must plan such use in advance
 - > “Planning” involves designing and implementing techniques that can efficiently manage these issues
 - > By working “off-line,” issues like naming become easier
 - > Most decisions can be made in small, constant amount of time

The Benefit

- Management & overhead time incurred once
 - Management & overhead time reduced by moving off line
- ⇒ The running program is faster

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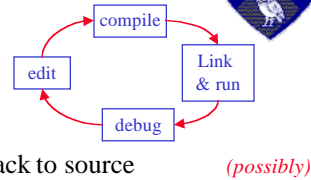
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The Complications of Using a Compiler



Requires the use of more tools

- Editor to enter program text
- Linker to tie compiled code together
- Debugger to relate post-execution state back to source



Environment is (typically) less standardized

- Depends on the libraries of code for “built-in” functionality
- Once we break the process up to allow compilation of parts, we can get into trouble with inconsistent or incomplete libraries
- The programmer now must keep track of these issues

Welcome to COMP 212 (and beyond)